## CMPE 350 - Spring 2017

## PS 11 - 09.05.17

- **4.30** Let A be a Turing-recognizable language consisting of descriptions of Turing machines,  $\{\langle M_1 \rangle, M_2 \rangle, \dots\}$ , where every  $M_i$  is a decider. Prove that some decidable language D is not decided by any decider  $M_i$  whose description appears in A. (Hint: You may find it helpful to consider an enumerator for A.)
- $\bullet$  Given an example of a language L such that L is co-Turing recognizable but its complement is not.
- Prove that the language  $\{< M, w, q > | M \text{ is a Turing machine which visits state } q \text{ during its execution when started with input string } w \}$  is undecidable.
- Show that the set of undecidable languages are closed under complementation.
- Some exam questions